Verified hybrid LQ control for autonomous spacecraft

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A benchmark problem for verified control

ARPOD problem: Automous rendezvous, proximity operation, and docking for spacecraft

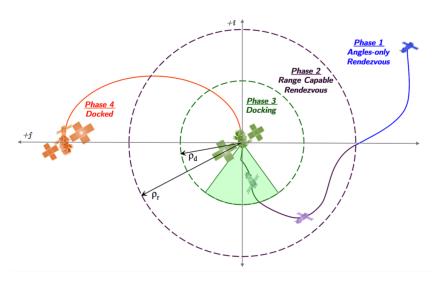
[Jewison and Erwin, CDC 2016]

Hybrid dynamical system

Control design

Automatic safety verification





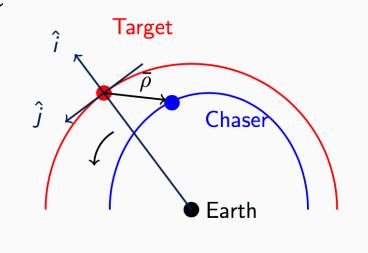
Plant model

State vector: $\bar{x} = [x, y, \dot{x}, \dot{y}]$

Input vector: $\bar{u} = [F_x, F_y]$

Separation: $\rho = \sqrt{x^2 + y^2}$

Angle of approach: $\theta = \arctan(\frac{y}{x})$



non-inertial relative coordinate frame with target located at the origin.

Underlying plant dynamics

Nonlinear

Derived from Kepler's laws and two-body problem

$$\ddot{x} = n^2 x + 2n\dot{y} + \frac{\mu}{r^2} - \frac{\mu}{r_c^3} (r + x) + \frac{F_x}{m_c}$$

$$\ddot{y} = n^2 y - 2n\dot{x} - \frac{\mu}{r_c^3} y + \frac{F_y}{m_c}$$

$$r_c = \sqrt{(r+x)^2 + y^2}$$
, $n = \sqrt{\frac{\mu}{r^3}}$, μ , r , m_c are given constants

Linear

Clohessy-Wiltshire-Hill (CWH) equations

$$\ddot{x} = n^2 x + 2n\dot{y} + \frac{F_x}{m_c}$$

$$\ddot{y} = -2n\dot{x} + \frac{F_y}{m_c}$$

$$\dot{\bar{x}} = A\bar{x} + B\bar{u} = (A - BK)\bar{x}$$

Modes in hybrid dynamics

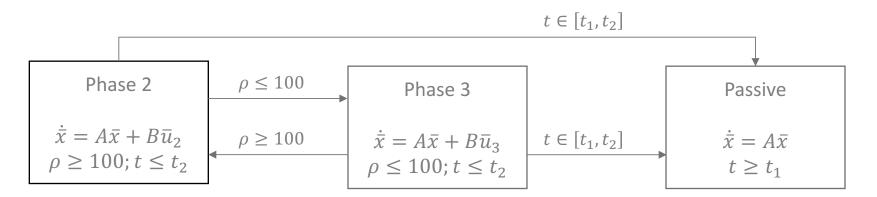
Phase 1: Only measure θ and not ρ ; system is not observable. (not in paper)

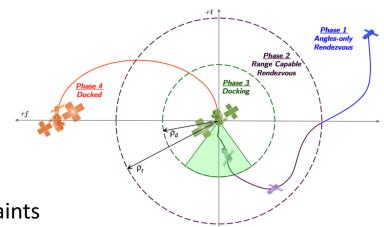
Phase 2: Chaser rendezvous with target without constraint.

Phase 3: Chaser continues rendezvous with constraints on its path and velocity, target location.

Phase 4: Plant mass changes and the terminal constraint is a new location. (not in this paper)

Abort/passive: Chaser shuts off its thrusters if a failure is detected





Safety constraints

Max thrust

$$|F_x|, |F_y| \le 10N$$

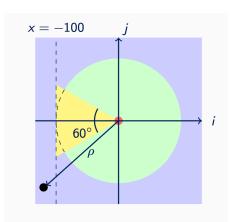
Closing velocity

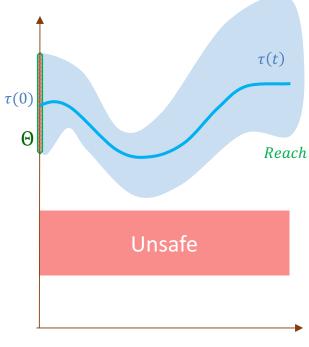
$$\sqrt{\dot{x}^2 + \dot{y}^2} \le 5cm/sec$$

Line of sight

$$\theta \in \left[-\frac{7\pi}{6}, -\frac{5\pi}{6}\right]$$

Problem: Design a controller for ARPOD and verify that all reachable states (from a set of initial states (Θ) and given set of disturbance inputs) meet these constraints.





Control and verification strategies

MPC-based controller using ellipsoidal constraints

[Jewison, Erwin, and Saenz-Otero 2015]

Optimal control using Reach-Avoid set computation

[Oishi et al. CDC 2016]

Hybrid supervisory control

[Malladi, Sanfelice, Butcher, and Wang, 2016-2017]

Trajectory planning using MPC (Phase 2) and differential flatness (Phase 1) [Farahani, Papusha, McGhan, and Murray]

Optimal control policy via stochastic reachability analysis

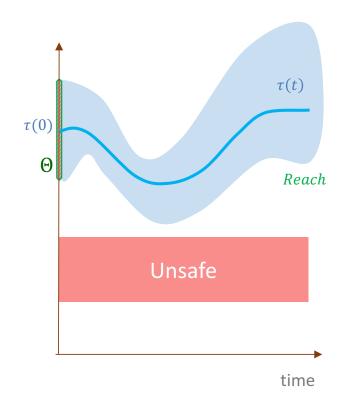
[Poonawala and Topcu, CDC 2016]

State-dependent LQ (SDLQ) and verification

[Chan and Mitra, 2016-17]

Simulation-driven bounded verification

Safety problem: given initial set Θ and unsafe set U, decide $Reach \cap U = \emptyset$?



Controller Design: SDLQ

$$\dot{\bar{x}} = A\bar{x} + B\bar{u} = (A - BK_i)\bar{x}$$

 Extend LQR to multiple stages; gives flexibility to handle local constraints

• Weights Q(·), R(·) of the quadratic cost are functions of the sampled state $\bar{x}(t)$; at ith period K_i is computed as:

$$\min_{\overline{u}} \int_0^\infty \overline{x}^T Q(\overline{x}(t_i)) \overline{x} + \overline{u}^T R(\overline{x}(t_i)) \overline{u}] dt$$

- Solution $K_i = R^{-1}B^TP_i$, where P_i is solution to algebraic Riccati equation
- Challenge: Simulations behave correctly, but analytical solution not available (needed for previous verification approaches)

Simulation-driven bounded verification

Simulation-driven verification for a single mode v

1. simulate \rightarrow 2. check safety \rightarrow 3. refine

Discrepancy β bounds distance between neighboring trajectories $\|\tau_1(t) - \tau_2(t)\| \le \beta(\tau_1(0), \tau_2(0), t)$,

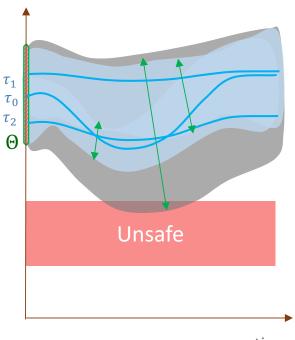
From a single simulation $\tau_1(t) + \beta$ over-approximate reach set from neighborhood of $\tau_1(0)$

Earlier approaches use f(x), $\frac{\partial f(x)}{\partial x}$

[C2E2: Duggirala et al. TACAS 15, Fan et al. CAV 15-16]

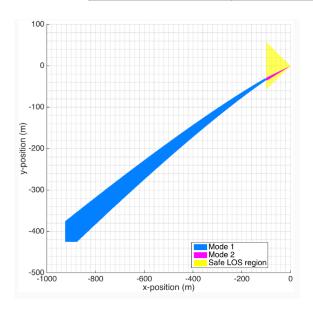
For LQR, closed-loop system admits analytical solution, sensitivity analysis, and verification with existing algorithms (tools like SpaceEx and C2E2)

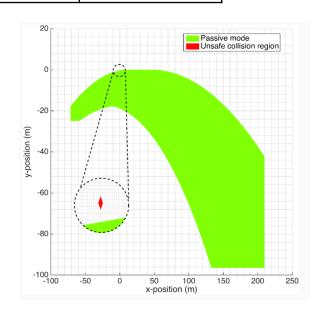
[Chan and Mitra, ARCH 2017]



Reachability analysis for LQ controller

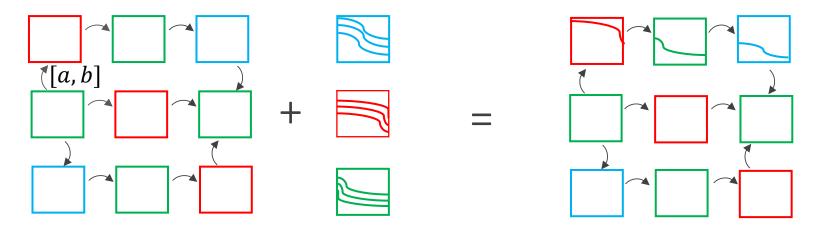
Algorithm	Linear with passive	Linear w/o passive	Nonlinear w/o passive
SDVTool [1]	Safe	Safe	n/v
SpaceEx [2]	Safe	Safe	n/v
C2E2 [3]	n/v	Safe	Safe





- [2] Frehse, et al. http://spaceex.imag.fr/
- [3] Duggirala, et al. http://publish.illinois.edu/c2e2-tool/
- [1] Chan and Mitra, MATLAB implementation of C2E2 algorithm

DryVR: A new view of hybrid verification



Transition graph Trace: $l_1, t_1, l_2, t_2, ..., l_k$ Black-box simulator

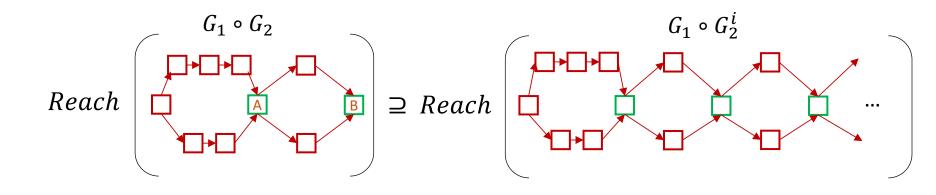
Trajectory: $\tau(t)$ Labeled trajectory set: $\langle \tau, l \rangle \in \mathcal{TL}$

Hybrid system $\mathcal{H} = \langle \mathcal{L}, \Theta, G, \mathcal{TL} \rangle$ State: a point in $\mathbb{R}^n \times \mathcal{L}$ $Reach = \{\langle x, l \rangle | \text{ for some } v, t, \langle x, l \rangle \text{ is }$ $reachable \text{ from } \Theta \}$ Reach|v: all states reachable in vertex v

[Fan, Qi, Mitra, and Viswanathan, CAV 2017]
[DryVR: http://dryvr.readthedocs.io/en/latest/index.html#]

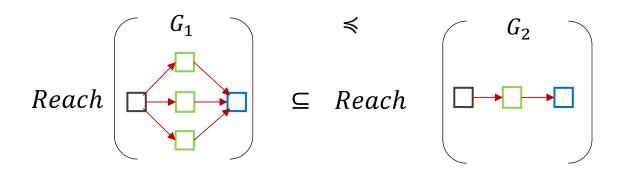
Composition for unbounded time analysis

If $Reach|B \subseteq Reach|A$ then



Reasoning about behavior containment

Trace containment $G_1 \leq G_2$ Trajectory containment $\mathcal{TL}_1 \leq \mathcal{TL}_2$ If $\Theta_1 \subseteq \Theta_2$, $G_1 \leq G_2$, $\mathcal{TL}_1 \leq \mathcal{TL}_2$, then



Learning discrepancy from data

Global exponential discrepancy function

$$\beta(x_1, x_2, t) = |x_1 - x_2| K e^{\gamma t}$$

For any pair of trajectories τ_1 and τ_2 in mode \square

$$\forall t \in [0, T], |\tau_1(t) - \tau_2(t)| \le |\tau_1(0) - \tau_2(0)| Ke^{\gamma t}$$

Taking logarithm and rearrange:

$$\forall t, \ln \frac{|\tau_1(t) - \tau_2(t)|}{|\tau_1(0) - \tau_2(0)|} \le \gamma t + \ln K$$

Learning linear separators

For $S \subseteq \mathbb{R} \times \mathbb{R}$, a linear separator is a pair $(a, b) \in \mathbb{R}^2$ s.t. $\forall (x, y) \in S, x \leq ay + b$ Algorithm:

- 1. Draw k pairs $(x_1, y_1), \dots, (x_k, y_k)$ from S according to \mathcal{D} .
- 2. Find $(a, b) \in \mathbb{R}^2$ such that $x_i \leq ay_i + b$ for all $i \in \{1, ..., k\}$.

Proposition [Valiant 84]: Let $\epsilon, \delta \in \mathbb{R}^+$. If $k \geq \frac{1}{\epsilon} \ln \frac{1}{\delta}$ then with probability $1 - \delta$, the above algorithm finds (a, b) such that $err_{\mathcal{D}}(a, b) < \epsilon$;

$$err_{\mathcal{D}}(a,b) = \mathcal{D}(\{(x,y) \in S \mid x > ay + b\})$$

Solve LP: min $2c \ln K + c(c+1)\gamma T$

s.t.
$$\forall i, j, s, \ln \frac{|\tau_i(t_s) - \tau_j(t_s)|}{|\tau_i(0) - \tau_j(0)|} \le \gamma t_s + \ln K$$

Bounded safety algorithm

Compute reach set from Θ : proceeds on G in a topologically sorted order

Refinement

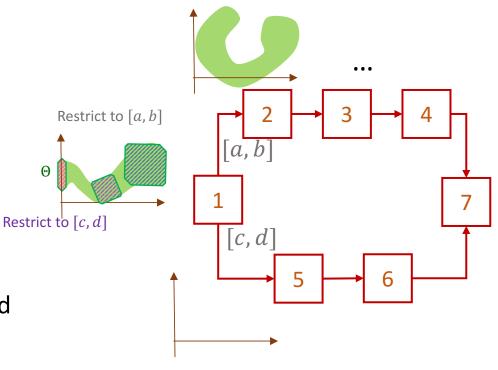
Split Θ to smaller sets

Split transition time intervals

Guarantee: Assuming that the learned discrepancy function is correct:

Soundness

Relative completeness

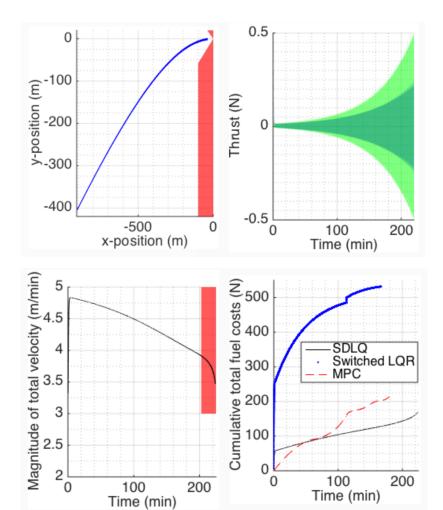


Conclusions

DryVR proves safety for Thrust and LOS constraints and a counterexample (unsafe) for Total Velocity constraint.

Simulation-driven verification, promising approach for grey-box models (try it)

Design and verification for complete ARPOD (with disturbance inputs)



Reachable positions (blue) and unsafe positions (red). (b) Reachable thrusts: Fx (blue) and Fy (green).

Composition for unbounded time analysis

If $Reach|B \subseteq Reach|A$ then

